

# **ENEE 459-C**

## **Computer Security**

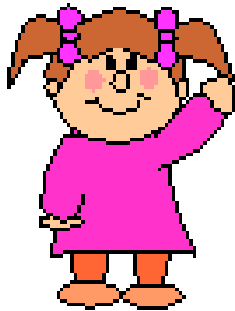
**Security protocols (continued)**



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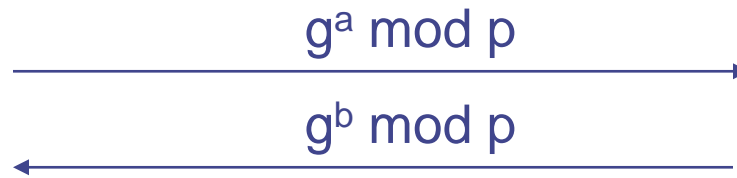
# Key Agreement: Diffie-Hellman Protocol

Key agreement protocol, both A and B contribute to the key  
Setup:  $p$  prime and  $g$  generator of  $Z_p^*$ ,  $p$  and  $g$  public.



Pick random, secret  $a$   
Compute and send  $g^a \bmod p$

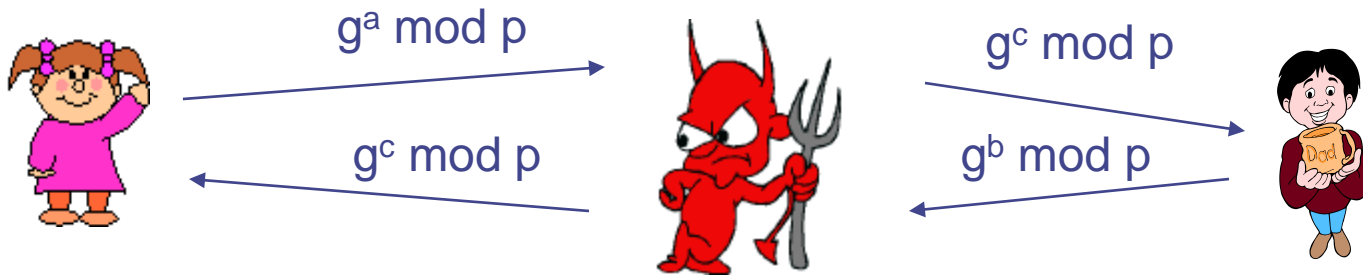
$$K = (g^b \bmod p)^a = g^{ab} \bmod p$$



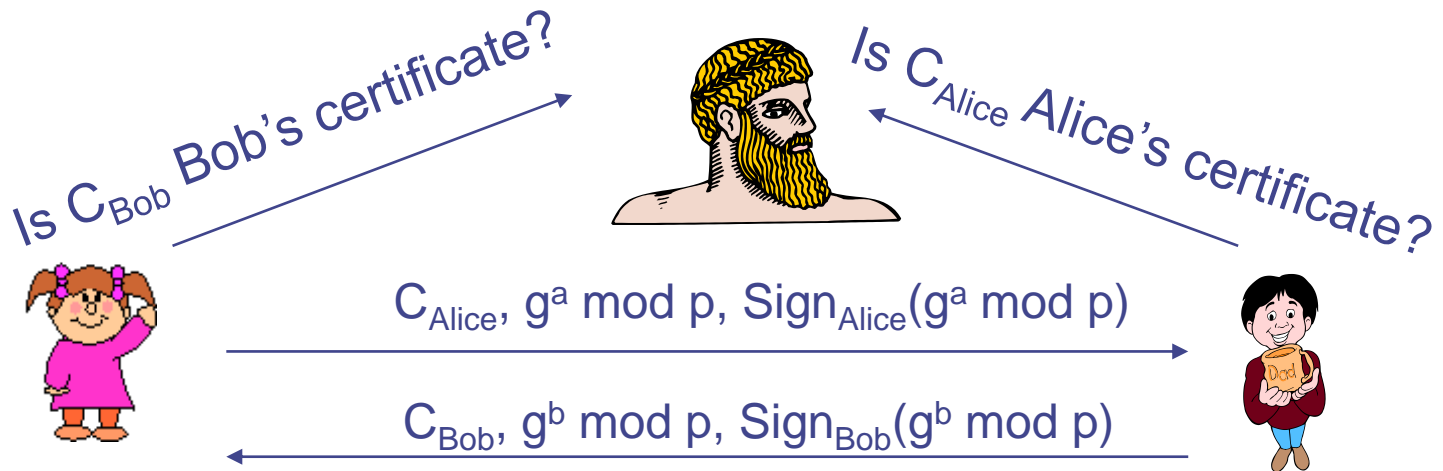
Pick random, secret  $b$   
Compute and send  $g^b \bmod p$

$$K = (g^a \bmod p)^b = g^{ab} \bmod p$$

# Authenticated Diffie-Hellman



Alice computes  $g^{ac} \bmod p$  and Bob computes  $g^{bc} \bmod p$  !!!



The slide features a light blue header bar at the top. A vertical blue line runs down the right side. Two horizontal blue lines intersect a vertical blue line on the left side, with a small circle at the intersection point. Another horizontal blue line intersects a vertical blue line on the right side, also with a small circle at the intersection point.

# Zero Knowledge

# Proofs of Knowledge

- An interactive proof system involves a prover and a verifier
- The prover proves he knows something to the verifier
- Zero-knowledge proof of knowledge (ZKPK):  
The prover convinces verifier that he knows a secret without revealing the secret
  - Application: Log into a server without revealing your password

# Properties of ZKPK

- Completeness

- If both prover and verifier are honest, protocol succeeds with overwhelming probability

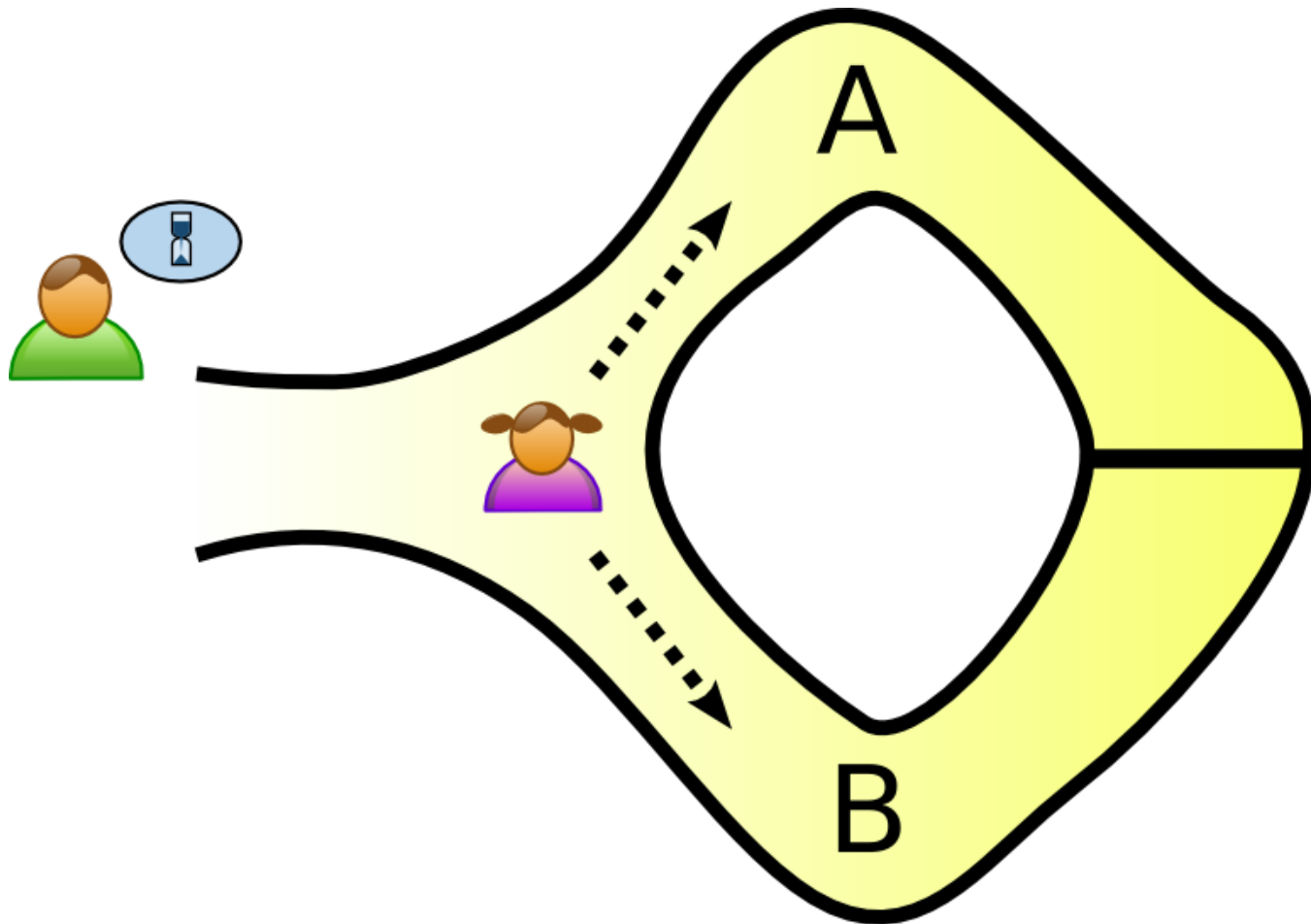
- Soundness

- No one who does not know the secret can convince the verifier (except with very small probability)
  - Intuition: the protocol should not enable prover to prove a false statement

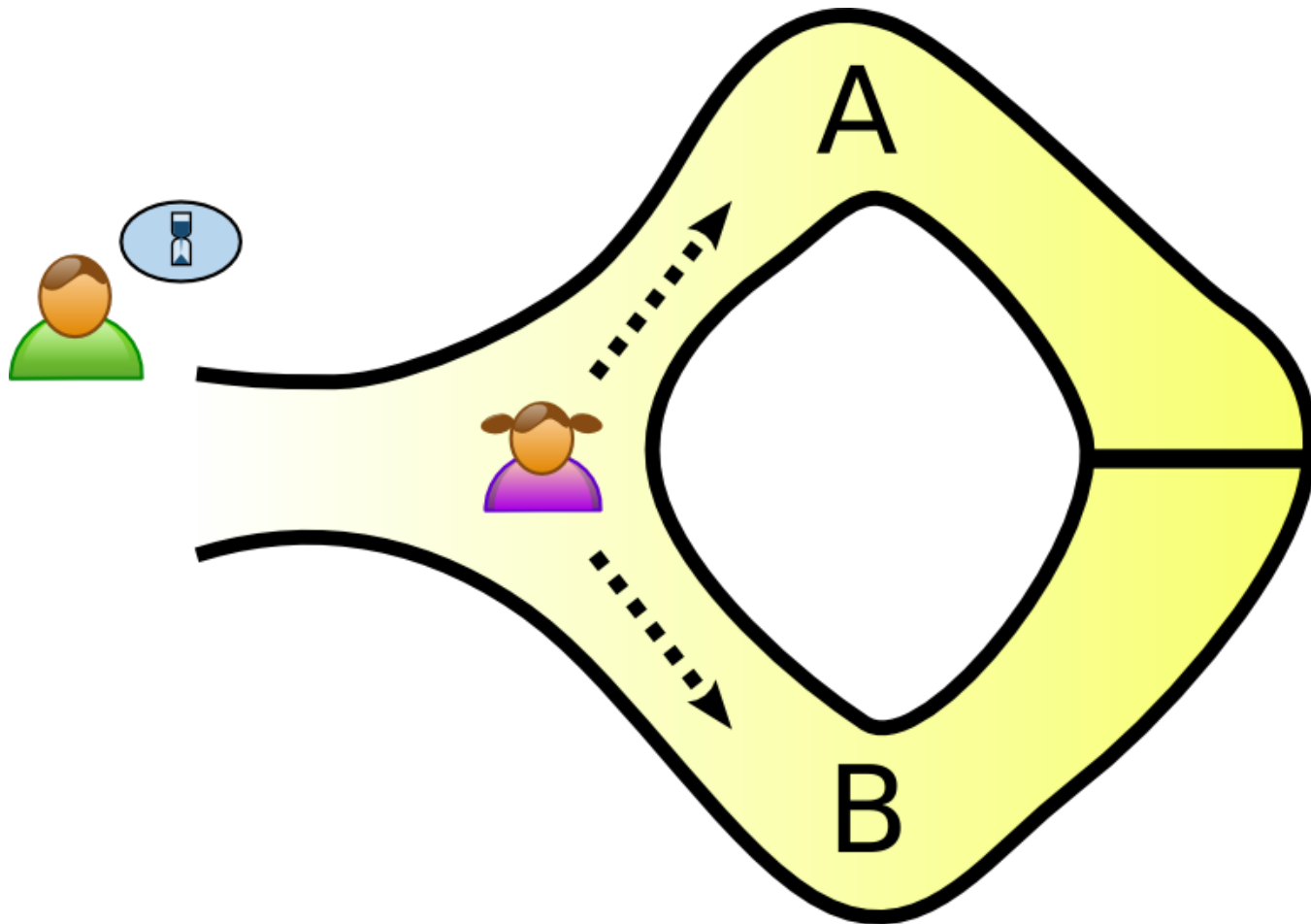
- Zero knowledge

- The proof does not leak any information

# An example (wikipedia)

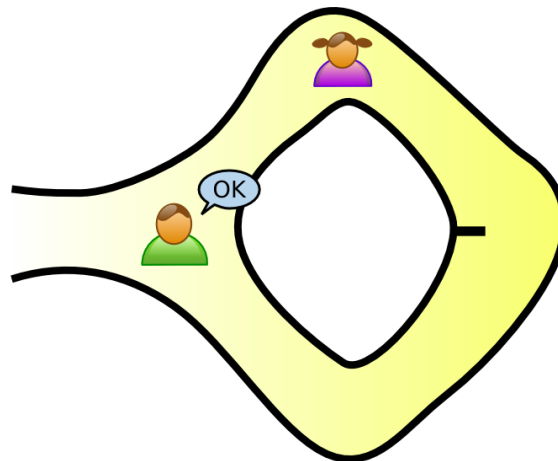
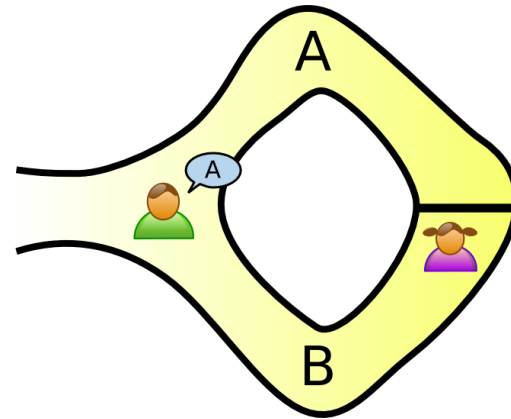
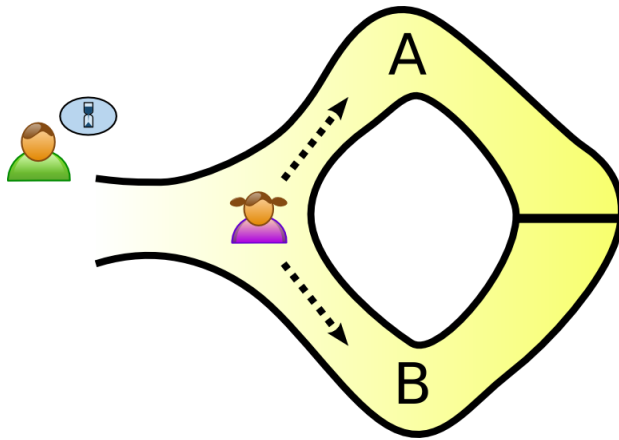


# An example (wikipedia)



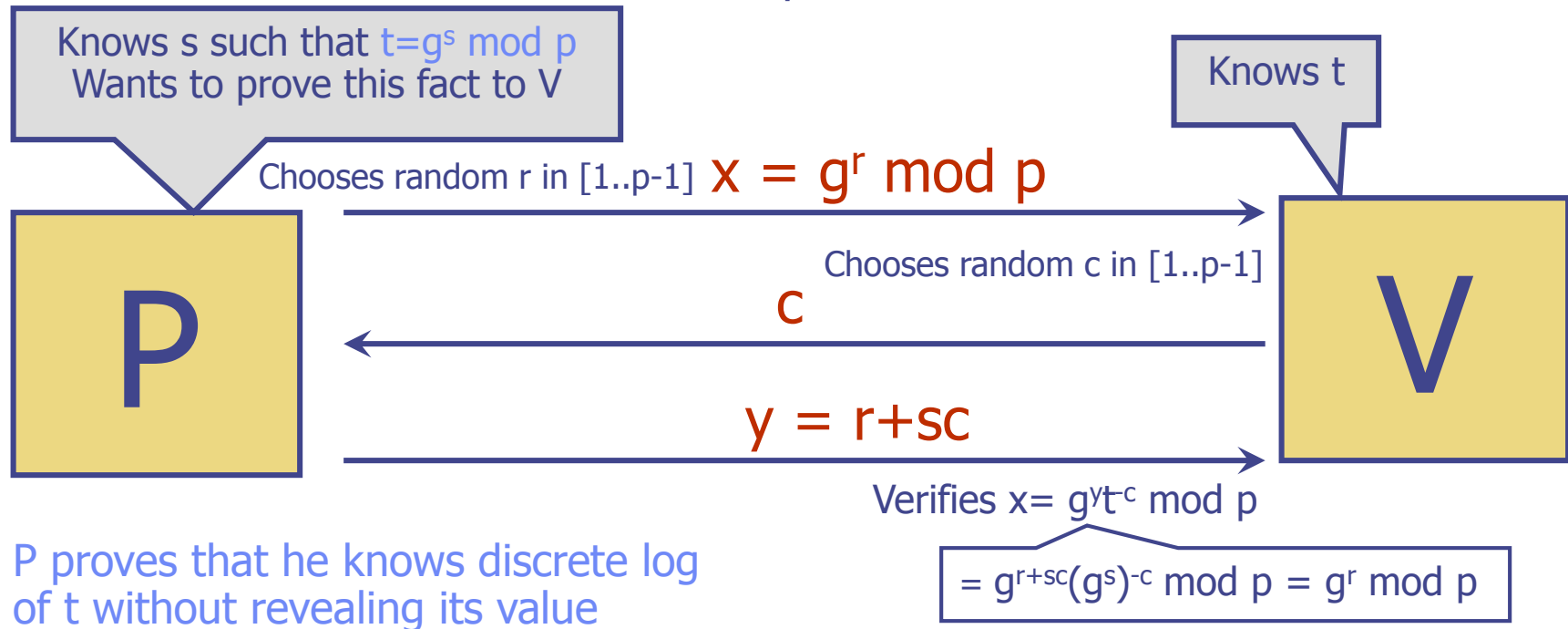


# An example (wikipedia)



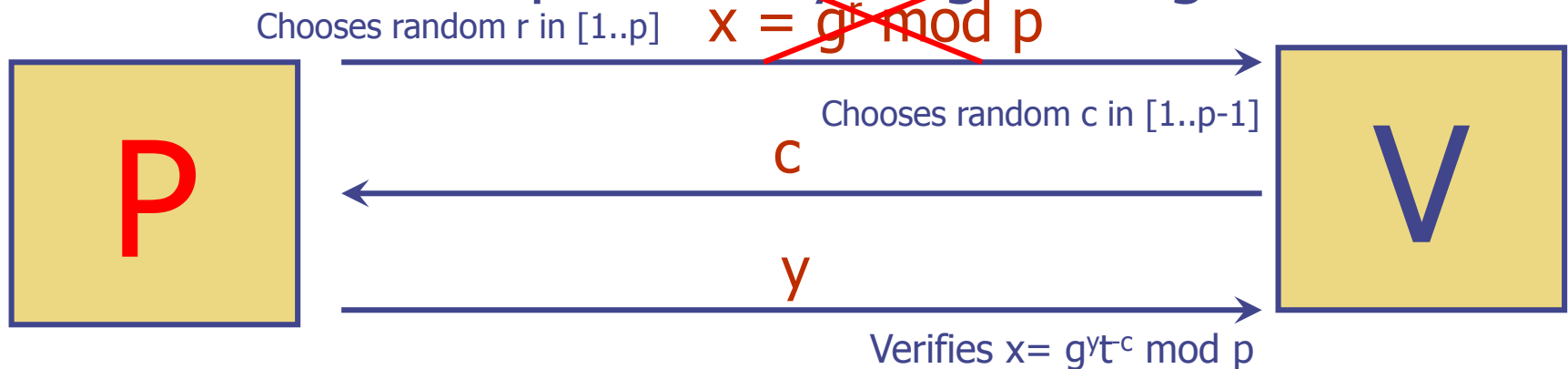
# Schnorr's Id Protocol

- System parameters
  - Prime  $p$
  - $g$  is a generator of  $\mathbb{Z}_p^*$



# Cheating Prover

- Prover can cheat if he can guess  $c$  in advance
  - Guess  $c$ , set  $x = g^y t^{-c}$  for random  $y$  in 1<sup>st</sup> message
  - What is the probability of guessing  $c$ ?



P proves that he "knows" discrete log of  $t$  even though he does not know  $s$

The slide features a minimalist design with thin blue lines. A vertical line runs down the left side, and a horizontal line runs across the top. Another vertical line is on the right, and a horizontal line is at the bottom. Small blue circular crop marks are located at the intersections of these lines: one at the top-left, one at the bottom-right, and one at the intersection of the left vertical line and the middle horizontal line. The title 'Anonymous Communication' is centered in a bold, orange, sans-serif font.

# Anonymous Communication

# Privacy on Public Networks

- Internet is designed as a public network
  - Machines on your LAN may see your traffic, network routers see all traffic that passes through them
- Routing information is public
  - IP packet headers identify source and destination
  - Even a passive observer can easily figure out **who is talking to whom**
- Encryption does not hide identities
  - Encryption hides payload, but not routing information

# What is Anonymity?

- Anonymity is the state of being not identifiable within a set of subjects
  - Hide your activities among others' similar activities
- Unlinkability of communicating identities
  - For example, sender and receiver are no more related after observing communication than they were before (difficult to achieve)

# Applications of Anonymity

- Privacy
  - Hide online transactions, Web browsing, etc. from intrusive governments, marketers and archivists
- Anonymous electronic voting
- Censorship-resistant publishing

# Attacks on Anonymity

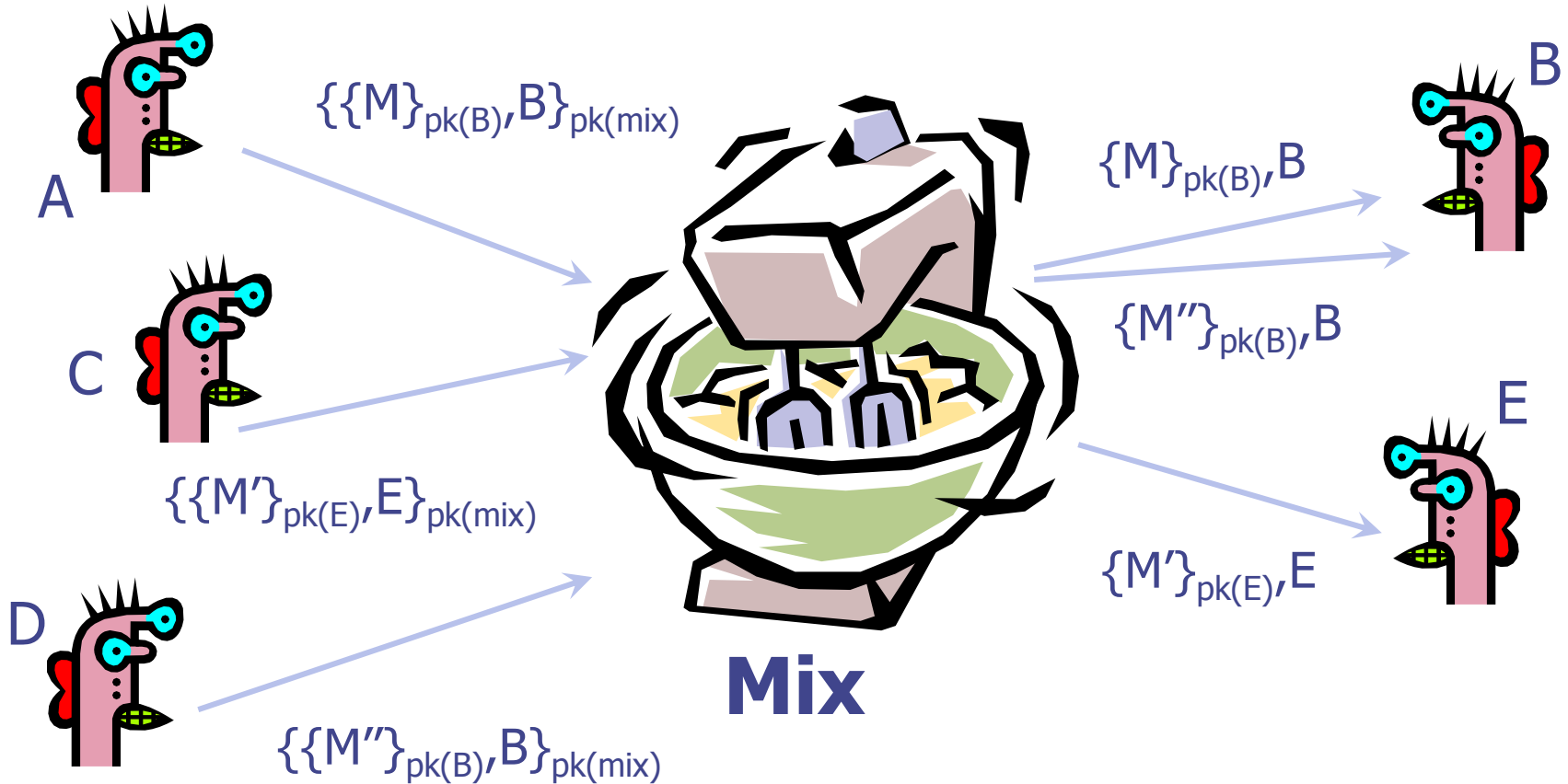
- Passive traffic analysis
  - Infer from network traffic who is talking to whom
- Compromise of network nodes
  - Attacker may compromise some routers
  - It is not obvious which nodes have been compromised
    - Attacker may be passively logging traffic
  - Better not to trust any individual router
    - Assume that some fraction of routers is good, don't know which



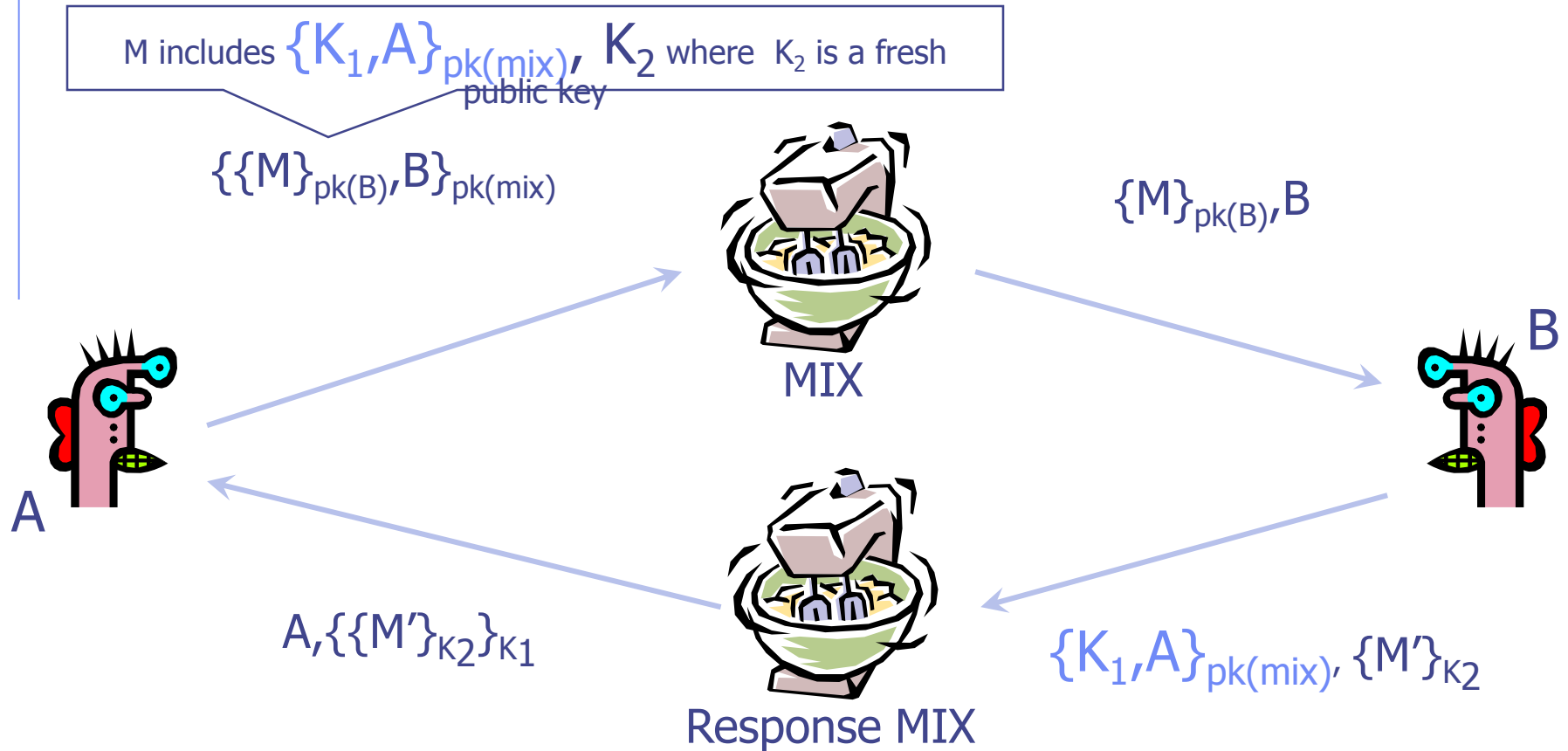
# Chaum's Mix

- Early proposal for anonymous email
  - David Chaum. "Untraceable electronic mail, return addresses, and digital pseudonyms". Communications of the ACM, February 1981.
- Public key crypto + trusted re-mailer (Mix)
- Modern anonymity systems use Mix as the basic building block

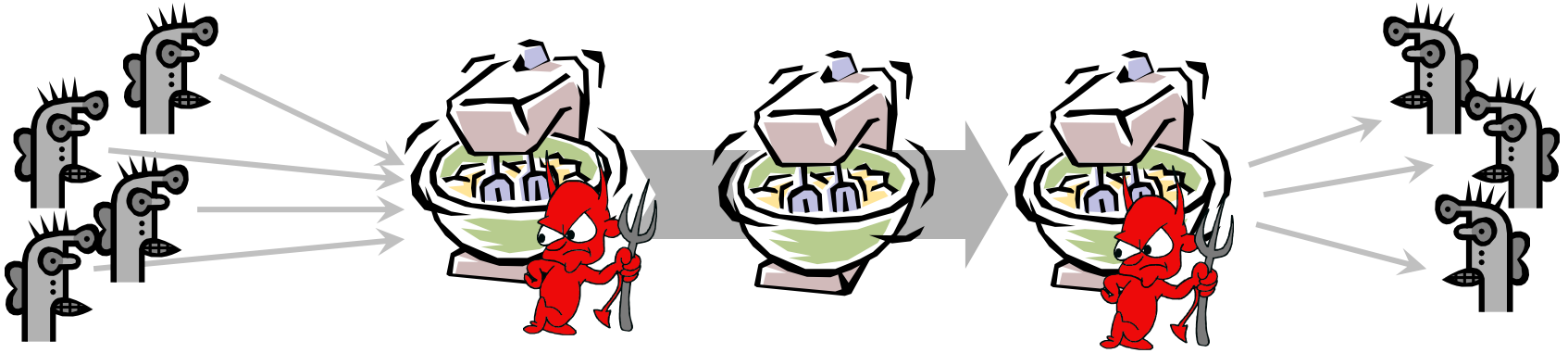
# Basic Mix Design



# Anonymous Return Addresses

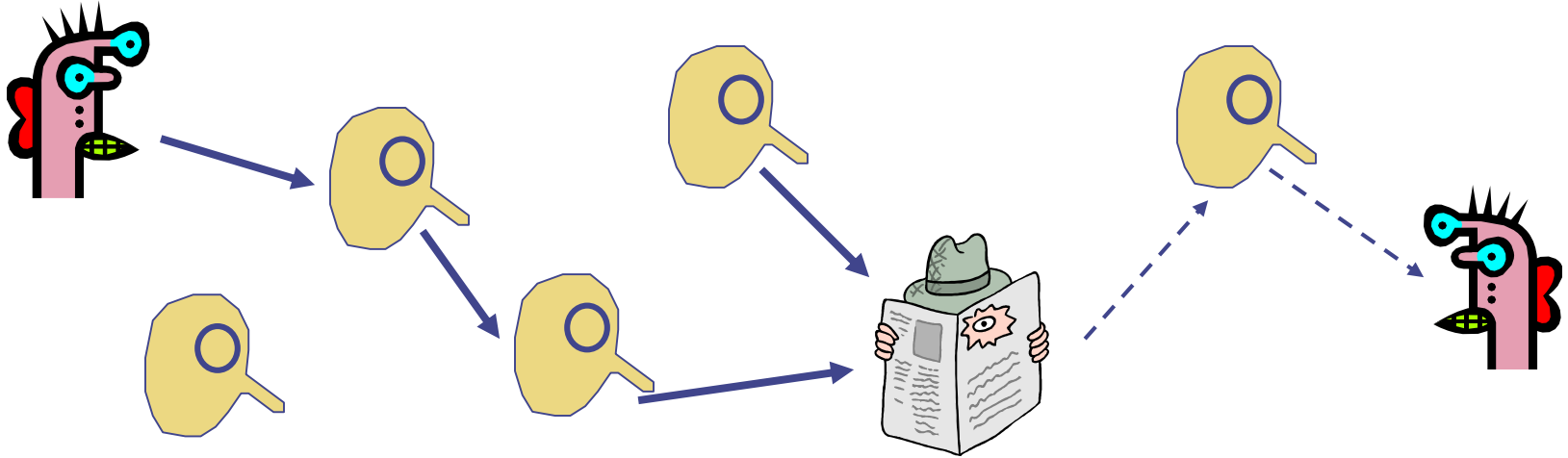


# Use many mixers



- Messages are sent through a **sequence of mixes**
  - Can also form an arbitrary network of mixes ("mixnet")
- Some of the mixes may be controlled by attacker, but even a single good mix guarantees anonymity

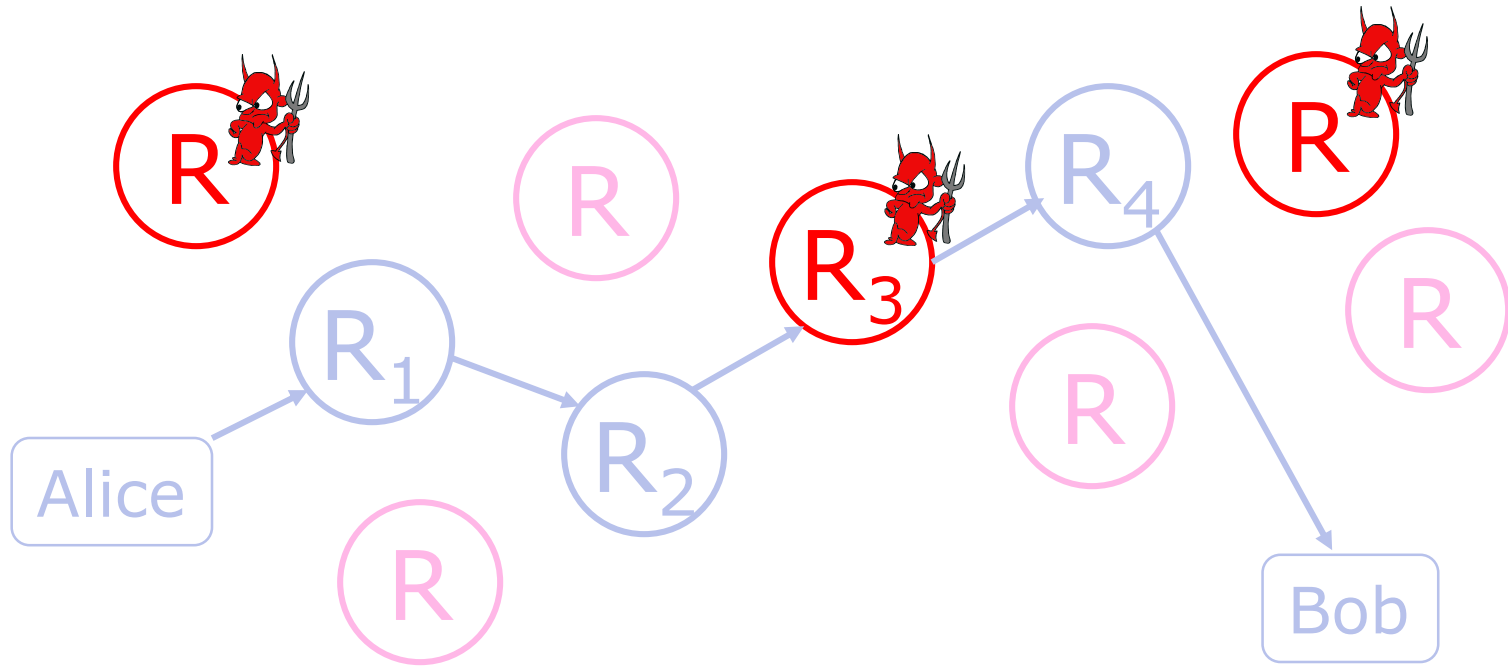
# Idea: Randomized Routing



- Hide message source by routing it randomly
  - Popular technique: Crowds, Freenet, Onion routing
- Routers don't know for sure if the apparent source of a message is the true sender or not

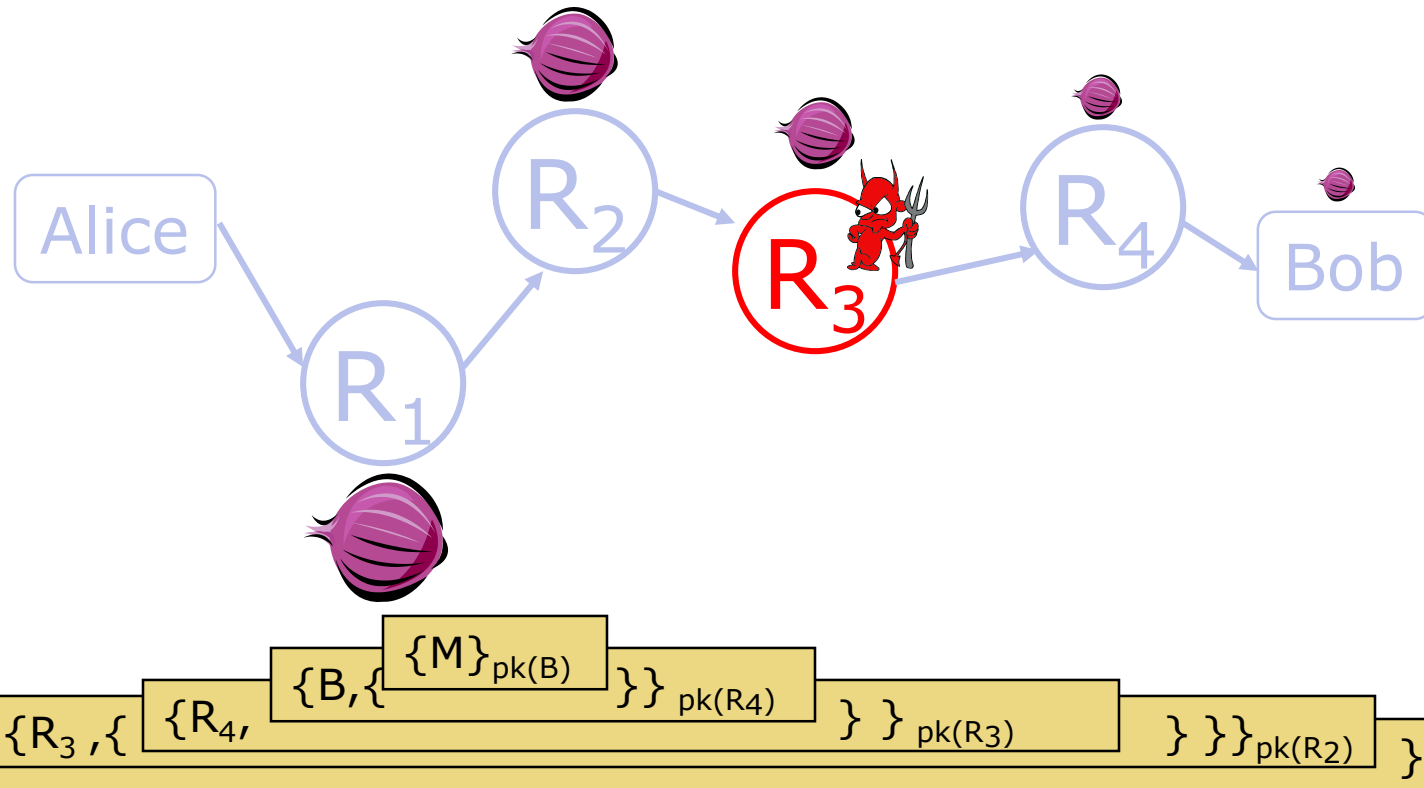
# Onion Routing

[Reed, Syverson, Goldschlag '97]



- ▶ Sender chooses a random sequence of routers
  - ▶ Some routers are honest, some controlled by attacker
  - ▶ Sender controls the length of the path

# Route Establishment



- Routing info for each link encrypted with router's public key
- Each router learns only the identity of the next router

# Disadvantages of Basic Mixnets/Onion Routing

- Public-key encryption and decryption at each mix/router are computationally expensive
- Basic mixnets have high latency
  - Ok for email, not Ok for anonymous Web browsing
- Challenge: low-latency anonymity network

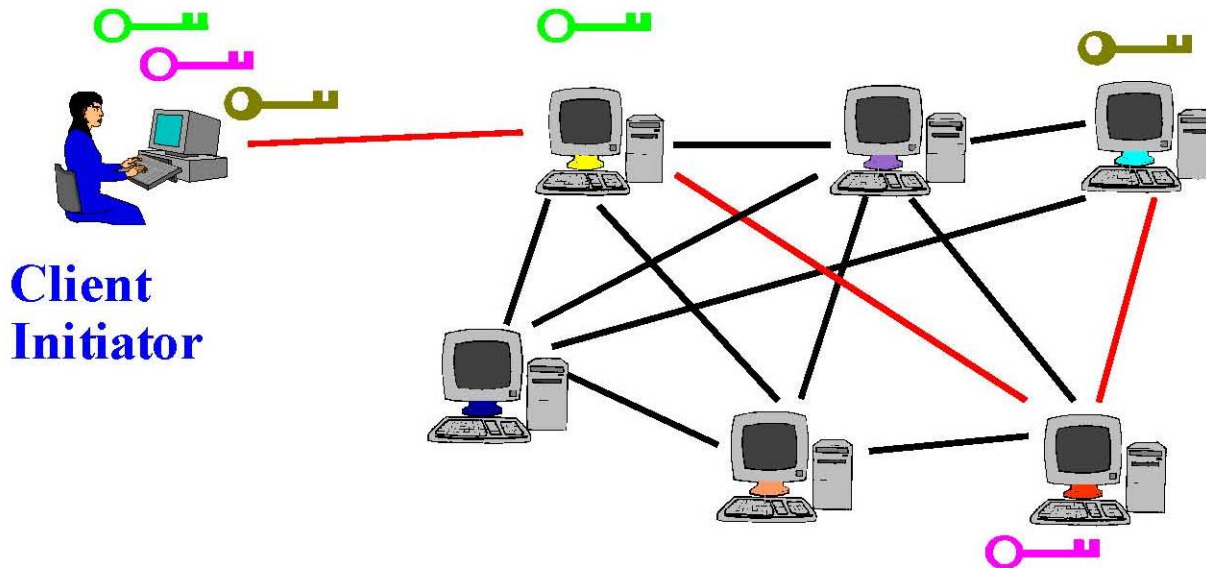


# Tor

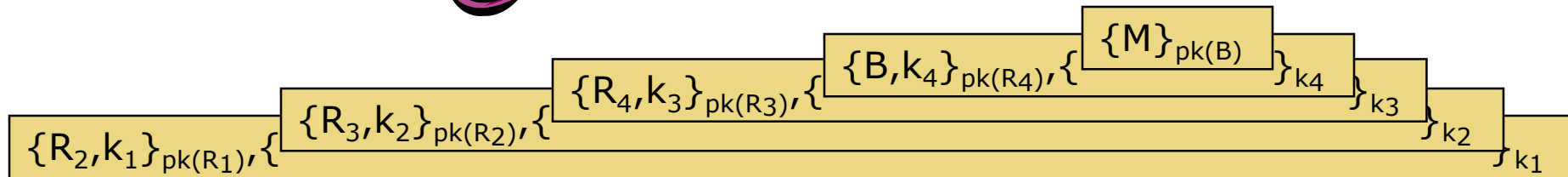
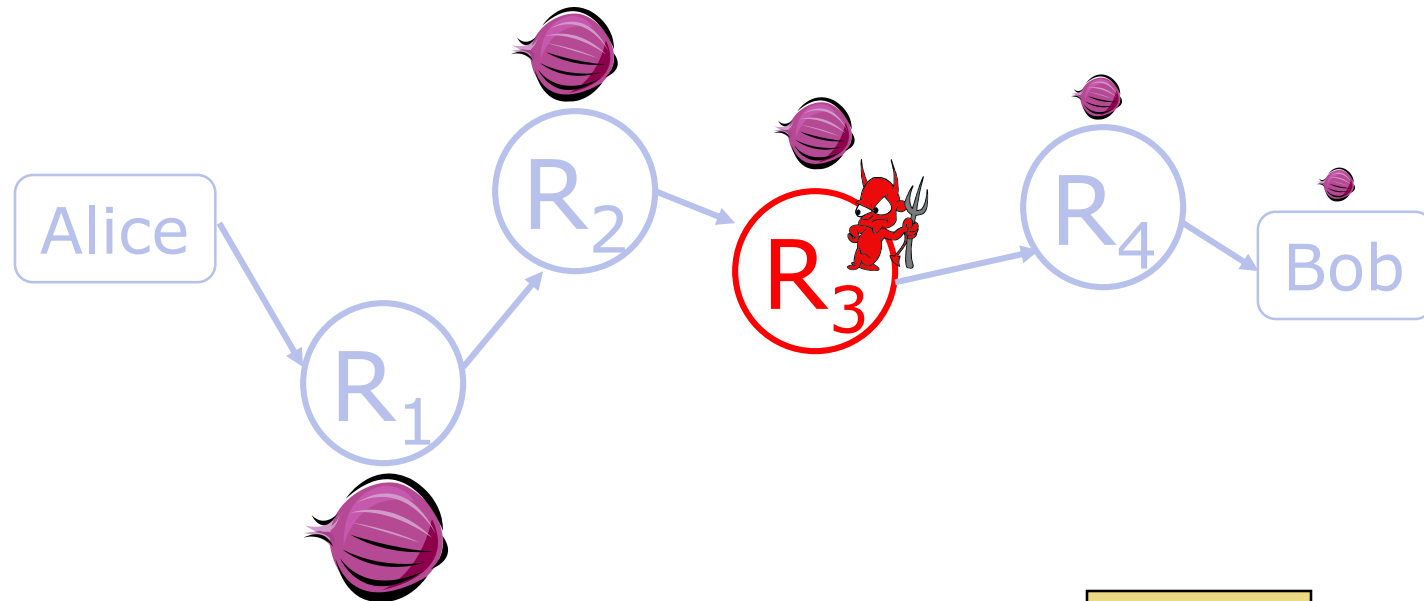
- Second-generation onion routing network
  - <https://www.torproject.org/>
  - Developed by Roger Dingledine, Nick Mathewson and Paul Syverson
  - Specifically designed for low-latency anonymous Internet communications
- Running since October 2003
- Thousands of users
- “Easy-to-use” client proxy
  - Freely available, can use it for anonymous browsing

# Tor Circuit Setup

- Client proxy establishes symmetric session keys with onion routers



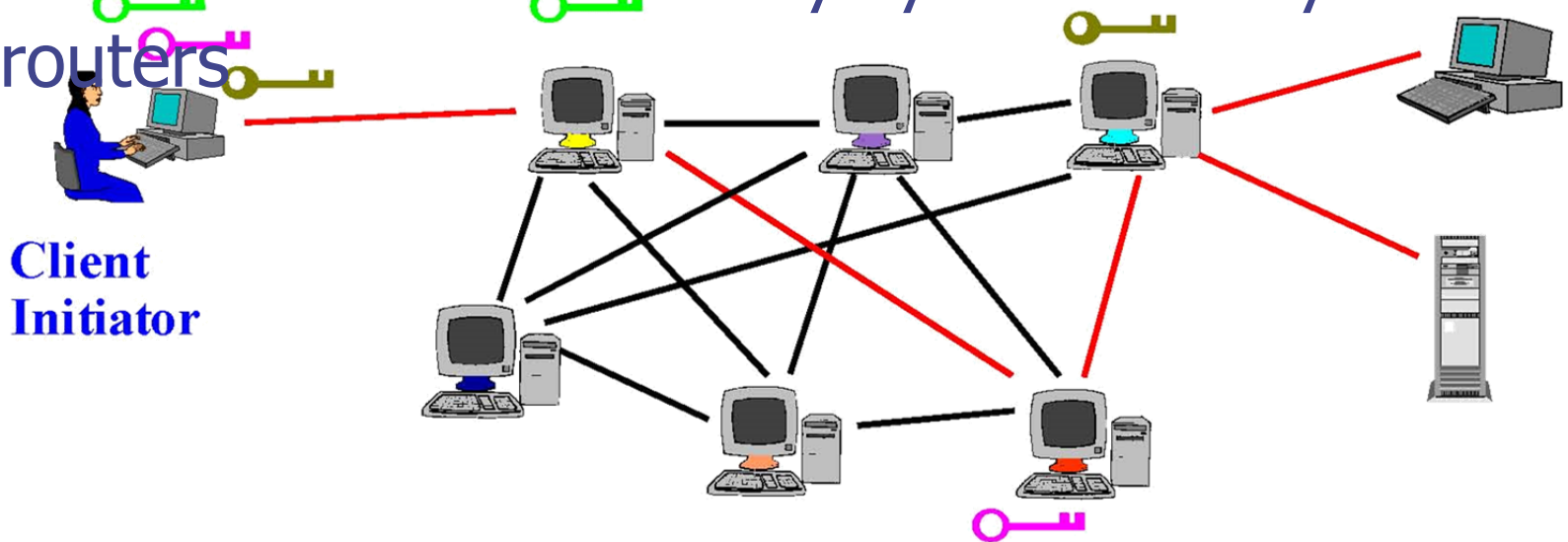
# Tor Circuit Setup (details)



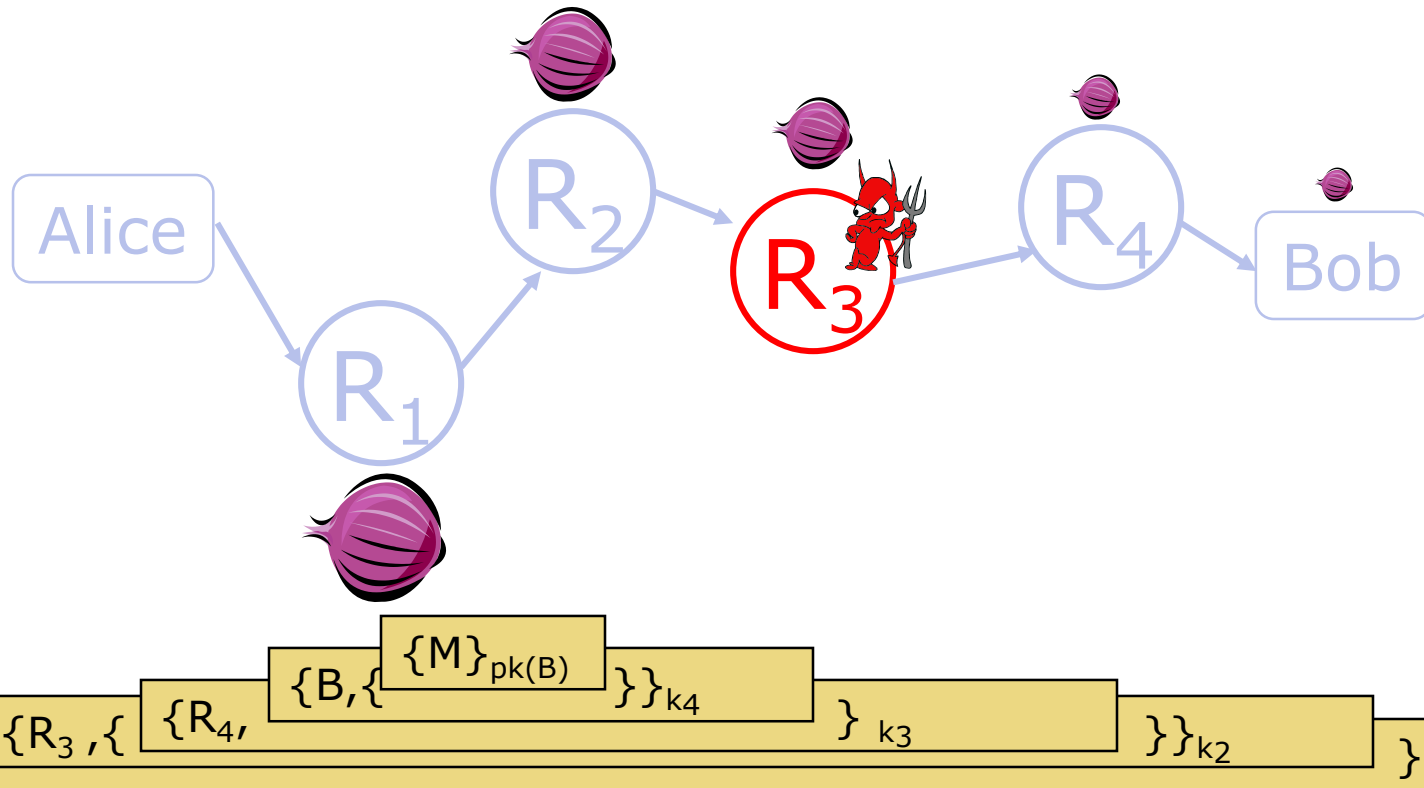
- Routing info for each link encrypted with router's public key
- Each router learns only the identity of the next router *and symmetric key with source*

# Using a Tor Circuit

- Client applications connect and communicate over the established Tor circuit
- Note onion now uses only symmetric keys for routers



# Using a Tor Circuit(details)



Note onion now uses only symmetric keys for routers

# Tor Management Issues

- Many applications can share one circuit
  - Multiple TCP streams over one anonymous connection
- Tor router doesn't need special privileges
  - Encourages people to set up their own routers
  - More participants = better anonymity for everyone
- Directory servers
  - Maintain lists of active onion routers, their locations, current public keys, etc.
  - Control how new routers join the network
    - "Sybil attack": attacker creates a large number of routers
  - Directory servers' keys ship with Tor code

# Deployed Anonymity Systems

- Tor (<http://tor.eff.org>)
  - Overlay circuit-based anonymity network
  - Best for low-latency applications such as anonymous Web browsing
- Mixminion (<http://www.mixminion.net>)
  - Network of mixes
  - Best for high-latency applications such as anonymous email